# Quantified cost functions for reasoning under uncertainty

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Cédric Pralet, Thomas Schiex, Gérard Verfaillie Mathematics of CS: Algebra, Logic and Graph Theory

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## Decision problems, $\operatorname{SAT}$ and $\operatorname{CSP}$

#### Constraint network

- a set X = {x<sub>1</sub>,..., x<sub>n</sub>} of variables with boolean/finite domains.
- a set C = {c<sub>1</sub>,..., c<sub>n</sub>} of boolean functions (clauses, constraints) with a scope (subset of X).

#### Query

Satisfiability:  $\exists x_1 \dots x_n (c_1 \wedge \dots \wedge c_n)$ 

# Optimisation: MaxSAT/CSP, Valued/Semiring CSP

### Cost function network

- a set X = {x<sub>1</sub>,..., x<sub>n</sub>} of variables with boolean/finite domains.
- a set C = {c<sub>1</sub>,..., c<sub>n</sub>} of cost functions over a set E (combination operator ⊗) with a scope.

#### Query

Optimisation:  $\max x_1 \dots \max x_n (c_1 \otimes \dots \otimes c_n)$ .

Addition: an utility set E, a combination op.  $\otimes$ , an elimination op. max, axioms.

#### Example

John has three doors (A, B, C) in front of him. Behind one, there is a treasure, and behind another a gangster. John may choose one door to open. He must pay  $4 \in$  to the gangster and the treasure amounts to  $10 \in$ .

#### The network

• Variables: ga, 
$$tr \in \{A, B, C\}$$
 and  $do \in \{A, B, C, nd\}$ .

• Cost functions:  $U_1 : do = ga(-4), U_2 : do = tr(10).$ 

No uncertainty: ga = A, tr = C. do = C (we get  $10 \in$ ).

# Adding uncertainty (plausibility)

### Non determinism

The treasure and the gangster are not behind the same door. Extra constraint:  $P_1 : ga \neq tr$ . Pessimistic query:  $\exists do \forall ga, tr, (P_1) \rightarrow (do \neq ga \land do = tr)$ ?

### Quantitative (probabilistic) uncertainty

Situations are all as likely. Cost function:  $P_2: \frac{1}{6}$  (normalization). Query: maximize expected utility

$$\max_{do} \sum_{tr,ga} \left( \left( \prod_{i=1}^{2} P_i \right) \times \left( \sum_{i=1}^{2} U_i \right) \right)$$

Two types of cost functions (P, U), specific combination operators, an operator to combine P and U. Multiple elimination (quantification) operators: max (decision), + (environment), =

#### Observations

John can listen to one door and try to detect the gangster. The probability of hearing is 0.8 if John listens to the ganster door but still 0.4 for a door next to this one. His friend Peter can do the same.

#### Network extension

Variables:  $Ii_J$ ,  $Ii_P$ , domain  $\{A, B, C\}$ ,  $he_J$ ,  $he_P$   $\{y, n\}$ Plausibilities:  $P_3 : P(he_J|ga, Ii_J)$ ,  $P_4 : P(he_P|ga, Ii_P)$  (norm)

**Query**: decision rules **maximizing the expected utility**, if first Peter and John listen in, and then John decides to open a door knowing what has been heared?

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$$\begin{bmatrix} \text{Solve the query} \\ \text{with a tree} \\ \text{exploration} \end{bmatrix} \Leftrightarrow \begin{bmatrix} \text{Record optimal decision rules for the quantity} \\ \max_{li_J, li_P} \sum_{he_J, he_P} \max_{do} \sum_{g^a, tr} \left(\prod_{i \in [1,4]} P_i\right) \times \left(\sum_{i \in [1,2]} U_i\right) \end{bmatrix}$$

### Other possible queries

### Opposite aims

If John thinks Peter is a traitor and let him choose a door to listen in first (**pessimistic attitude concerning the other agent**):

$$\min_{li_{P}} \max_{li_{J}} \sum_{he_{P}, he_{J}} \max_{do} \sum_{ga, tr} \left( \prod_{i \in [1, 4]} P_{i} \right) \times \left( \sum_{i \in [1, 2]} U_{i} \right)$$

#### Observabilities

If Peter does not even tell John what he has heard (John does not observe  $he_P$ ):

$$\min_{li_{P}} \max_{li_{J}} \sum_{he_{J}} \max_{do} \sum_{he_{P}} \sum_{ga,tr} \left( \prod_{i \in [1,4]} P_{i} \right) \times \left( \sum_{i \in [1,2]} U_{i} \right)$$

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## Adding feasibilities

John and Peter cannot listen in to the same door, door A is locked!



 

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John and Peter cannot listen in to the same door, door A is locked! Two local feasibility functions:  $F_1 : Ii_J \neq Ii_P$ ,  $F_2 : do \neq A$ 



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## Feasibility operators

Feasibilities is not unacceptibility. It restricts elimination domains.

• false 
$$\star \alpha = \Diamond$$
, true  $\star \alpha = \alpha$ 

•  $\Diamond$  is ignored by elimination operators (max( $\Diamond$ , u) = min( $\Diamond$ , u) =  $\Diamond$  + u = u)

Computation of **optimal decision rules** with:

$$\min_{li_{P}} \max_{li_{J}} \sum_{he_{J}} \max_{do} \sum_{he_{P}} \sum_{ga,tr} \left( \left( \bigwedge_{i \in [1,2]} F_{i} \right) \star \left( \prod_{i \in [1,4]} P_{i} \right) \times \left( \sum_{i \in [1,2]} U_{i} \right) \right)$$

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## Many frameworks, many related algorithms



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### Three key components

- the algebraic structure, defining properties satisfied by the combination and the elimination operators
- the network of local functions, with normalization conditions on plausibilities and feasibilities
- the **query** (sequence of eliminations).

## An algebraic structure for plausibilities, utilities...

Extension of Halpern and Chu work to sequential decisions.

- plausibilities:  $(E_p, \leq_p, \oplus_p, \otimes_p)$  commutative semiring with monotonic operators.
- **2 utilities**:  $(E_u, \preceq_u, \otimes_u)$  commutative monoid.
- **3** expected utility:  $(E_p, E_u, \oplus_u, \otimes_{pu})$   $(E_u, \oplus_u, \otimes_{pu})$  semi-module on plausibilities:
  - $\otimes_{pu}$  distributive wrt  $\oplus_p$  and  $\oplus_u$ .
  - $\forall p_1, p_2 \in E_p, \forall u \in E_u, p_1 \otimes_{pu} (p_2 \otimes_{pu} u) = (p_1 \otimes_p p_2) \otimes_{pu} u.$
  - $\forall u \in E_u, (0_p \otimes_{pu} u = 0_u) \land (1_p \otimes_{pu} u = u).$

$$\sum_{j} (p_j \times u_j) \text{ becomes } \bigoplus_{j} (p_j \otimes_{pu} u_j)$$

	Ep	$\oplus_{p}$	$\otimes_{p}$	Eu	$\otimes_{u}$	$\oplus_{u}$	$\otimes_{\it pu}$
1	$\mathbb{R}^+$	+	Х	$\mathbb{R} \cup \{-\infty\}$	+	+	×
2	$\mathbb{R}^+$	+	Х	$\mathbb{R}^+$	×	+	×
3	[0, 1]	max	min	[0, 1]	min	max	min
4	[0,1]	max	min	[0,1]	min	min	$\max(1-p, u)$
5	$\mathbb{N}\cup\{\infty\}$	min	+	$\mathbb{N}\cup\{\infty\}$	+	min	+
6	$\{t, f\}$	$\vee$	$\wedge$	$\{t, f\}$	$\wedge$	V	$\wedge$
7	$\{t, f\}$	V	$\wedge$	$\{t, f\}$	$\wedge$	$\wedge$	$\rightarrow$

(1) probabilistic expected utility, (2) probabilistic expected satisfaction, (3/4) optimistic/pessimistic possibilistic utilities, (5)  $\kappa$ -rankings with positive utilities (6/7) optimistic/pessimistic expected boolean satisfaction

## Locality of functions

### **Definition**: a **PFU network** is a tuple (V, G, P, F, U) where:

- V: finite set of variables, partitionned between V<sub>D</sub> the set of decision variables, and V<sub>E</sub>, the set of environment variables. Each is partitioned in clusters.
- *G* a a cluster DAG representing normalization/independance.
- $P = \{P_1, P_2, \ldots\}$  a finite set of local plausibility functions
- $F = \{F_1, F_2, \ldots\}$  a finite set of local feasibility functions
- $U = \{U_1, U_2, \ldots\}$  a finite set of local utility functions

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## Example



### Query Q on a PFU network $\mathcal{N}$

- A pair  $(\mathcal{N}, Sov)$  where:
  - $\mathcal{N}$ : PFU network
  - Sov is a sequence of operator/variable(s) pairs. Operators are min, max, or  $\bigoplus_u$  and act as quantifiers.

### $Ex.: (\mathcal{N}, \max_{x_1, x_2} \min_{x_3} \oplus_{u \times 4, x_5} \max_{x_6})$ Meaning:

- the order specifies the order in which decisions are made and environment variables are observed
- min and max operators applied on decision variables specify if an optimistic or a pessimistic attitude is adopted for decisions.

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## Correct queries

- Adequation between an eliminated variable and its elimination operator (decision variable → eliminated with min or max; environment variables → eliminated with ⊕<sub>u</sub>). Ex: if Sov = ... ∑<sub>do</sub>..., the query is not correct
- Respect of causality: conditions (imposed by the DAG) on the order of elimination
  - Ex: if  $Sov = \ldots \sum_{he_J} \ldots \max_{li_J} \ldots$ , the query is not correct.

# Value of a query

### Answer to a query Q

The answer Ans(Q) to a correct query Q can be defined based on **decision trees**.

- Advantage: has a clear semantic foundation.
- Drawback: weights associated with some edges in the tree must be computed, which can take an exponential time.

#### Property

The definition is **equivalent** to the following operational one:

$$Ans(Q) = Sov\left(\left(\bigwedge_{F_i \in F} F_i\right) \star \left(\bigotimes_{P_i \in P} P_i\right) \otimes_{pu} \left(\bigotimes_{U_i \in U} U_i\right)\right)$$

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**Theorem**: the PFU framework can be used to model and solve queries asked on:

CSP Valued CSP Quantified CSP Mixed and probabilistic CSP Stochastic CSP

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SAT Quantified boolean formulas Stochastic SAT Extended stochastic SAT

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## Not covered

- non distributional uncertainty (Dempster-Shafer)
- really qualitative utility (CP-nets) or partially ordered utilities (Queries).

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The operational definition gives us a **generic tree search algorithm** (and PSPACE membership).

Variable elimination: additional axioms needed (sufficient).

- (Ax1) semiring:  $((E_p, \oplus_p, \otimes_p) = (E_u, \oplus_u, \otimes_u)) \land (\otimes_{pu} = \otimes_u) \land (\preceq_p = \preceq_u)$ (2,3,5,6)
- **2** (Ax2) vector space-like:  $\oplus_u = \otimes_u$  on  $E_u$  (1,4,7).



- combine existing algorithms: branch and bound, local consistency (csp, hard information), cost function consistency (soft constraints), tree decomposition based (recursive conditioning).
- variable elimination/CTE: constrained treewidth and quantified cost functions: connection between the query and the network structure. Do query optimization to lower contrained treewidth.
- islands of tractability/NP-completeness and beyond (multi-chotomy ?) in quantified cost-functions.