

Cost function optimization & local consistency

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Overview

☐ Introduction and definitions

- What is it ?
- How do we solve it ?

☐ Local consistency

- Fundamental operations on soft CN
- How are fair valuation structures ?
- Equivalence preserving operations
- Local consistencies
- Max flow as a local consistency ?

☐ Conclusion

Why soft constraints?

- CSP framework: for *decision* problems
- Many problems are *overconstrained* or *optimization* problems
- Economics (combinatorial auctions)
 - Given a set G of goods and a set B of bids...
 - Bid (B_i, V_i) , B_i requested goods, V_i value
 - ... find the best subset of compatible bids
 - Best = maximize revenue (sum)

Why soft constraints?

☐ Satellite scheduling

- Spot 5 is an earth observation satellite
- It has 3 on-board cameras
- Given a set of requested pictures (of different importance)...
 - ☐ Resources, data-bus bandwidth, setup-times, orbiting
- ...select a subset of compatible pictures with max. importance (sum)

Why soft constraints?

- Probabilistic inference (bayesian nets)
 - Given a probability distribution defined by a DAG of conditional probability tables
 - And some evidence
 - ...find the *most probable* explanation for the evidence (product)

Haplotyping

Why soft constraints?

- ☐ Ressource allocation (frequency assignment)
 - Given a telecommunication network
 - ...find the best frequency for each communication link avoiding interferences

- Best can be:
 - ☐ Minimize the maximum frequency (max)
 - ☐ Minimize the global interference (sum)

w/o mentionning MaxCut/Clique, Vertex cover...

Soft constraint network (CN)

□ (X, D, C)

■ $X = \{x_1, \dots, x_n\}$ variables

■ $D = \{D_1, \dots, D_n\}$ finite domains

■ $C = \{f, \dots\}$ e **cost functions**

□ $f_S, f_{ij}, f_i, f_\emptyset$ scope $S, \{x_i, x_j\}, \{x_i\}, \emptyset$

□ $f_S(t): \rightarrow E$ (ordered by $\preceq, \perp \preceq T$)

identity

annihilator

□ Obj. Function: $F(X) = \oplus f_S (X[S])$

□ Solution: $F(t) \neq T$

□ Soft CN: find **minimum cost** solution

- commutative
- associative
- monotonic

Specific frameworks

Instance	E	$\oplus$	$\perp \preceq T$
Classic CN	$\{t, f\}$	and	$t \preceq f$
Possibilistic	$[0, 1]$	max	$0 \preceq 1$
Fuzzy CN	$[0, 1]$	usual min	$1 \preceq 0$
Weighted CN (bounded or not)	$[0, k]$	$+$	$0 \preceq k$
Bayes net	$[0, 1]$	$\times$	$1 \preceq 0$

Lexicographic CN, probabilistic CN...

Basic operations

- combination / join
- projection

Cost function implication

f_P implied by g_Q ($f_P \preceq g_Q$, g_Q tighter than f_P)
 iff for all $t \in \ell(P \sqcap Q)$, $f_P(t[P]) \preceq g_Q(t[Q])$

x_i	x_j	$g(x_i, x_j)$
b	b	2
b	g	3
b	r	2
g	b	4
g	g	5
g	r	4
r	b	T=6
r	g	T=6
r	r	T=6

implies

x_i	$f(x_i)$
b	1
g	2
r	3

or

x_i	$g[x_i]$
b	2
g	4
r	T=6

T=1 : classic CSP case

Combination (join with $\oplus$, + here)

x_i	x_j	$f(x_i, x_j)$
b	b	6
b	g	0
g	b	0
g	g	6

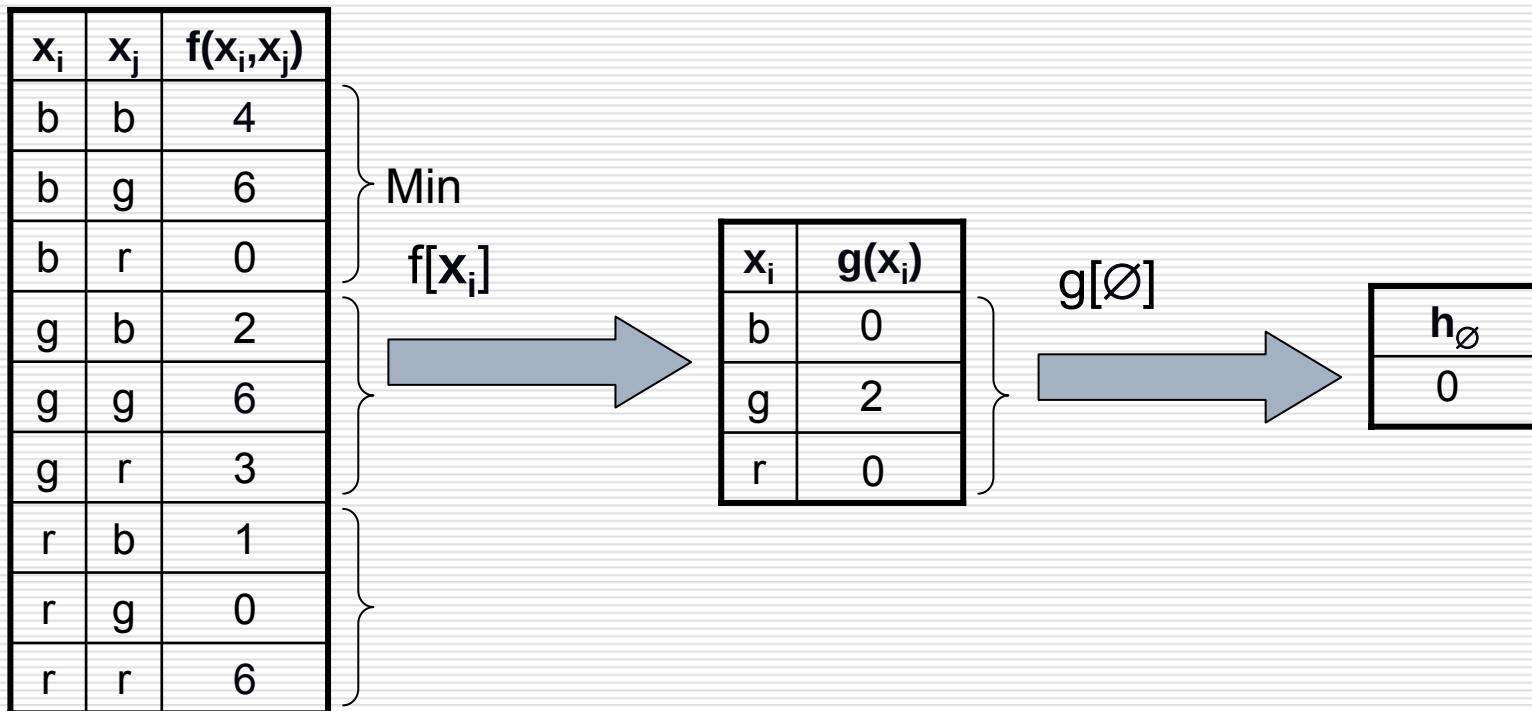
$\oplus$

x_i	x_j	x_k	$h(x_i, x_j, x_k)$
b	b	b	12
b	b	g	6
b	g	b	0
b	g	g	6
g	b	b	6
g	b	g	0
g	g	b	6
g	g	g	12

$= 0 \oplus 6$

x_j	x	$g(x_j, x_k)$
b	b	6
b	g	0
g	b	0
g	g	6

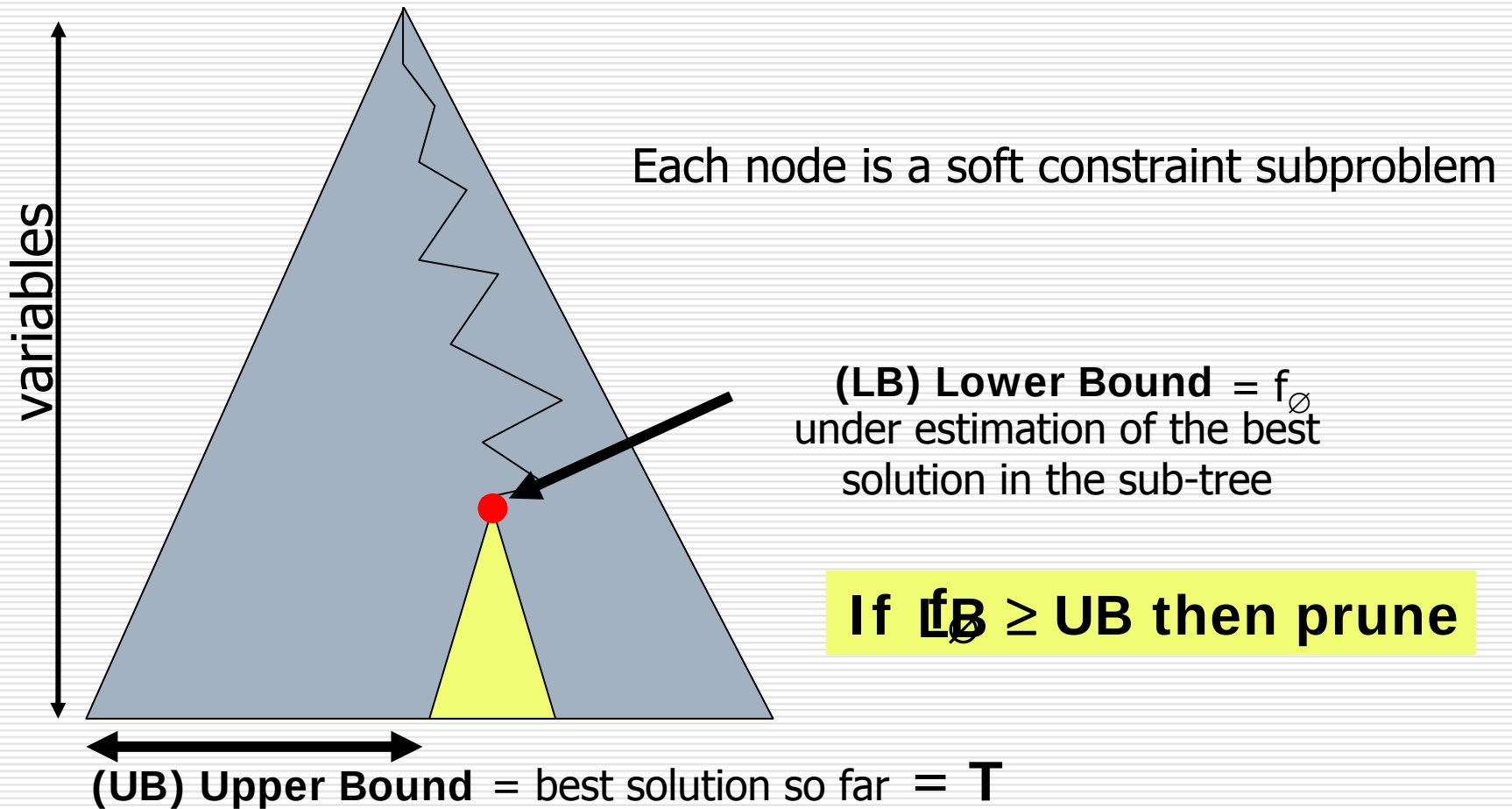
Projection (elimination)



Systematic search

Branch and bound(s)

Systematic search



Local consistency

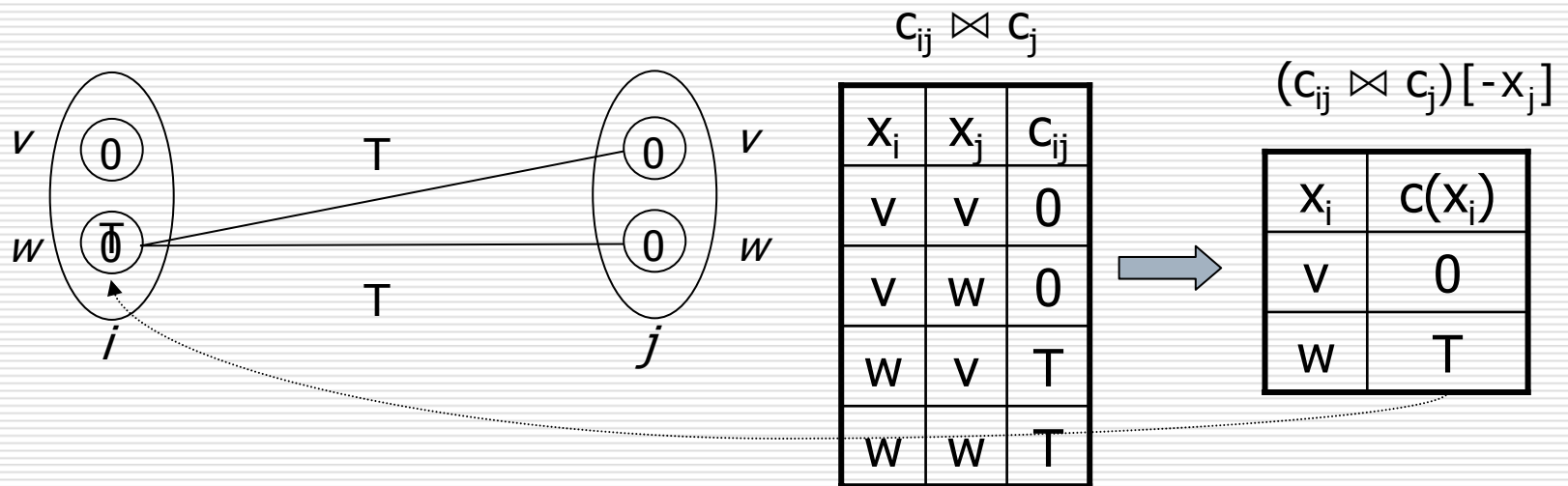
- Very useful in classical CSP
 - Node consistency (1-consistency)
 - Arc consistency...(2-consistency in binary CSP)
 - Related to tractable classes (AC: tree, 2SAT)

- Produces an equivalent more explicit problem:
 - Preserves solutions (costs), may detect unfeasibility (gives a better $f_{\emptyset}$)

Classical arc consistency (binary CSP)

□ for any x_i and c_{ij}

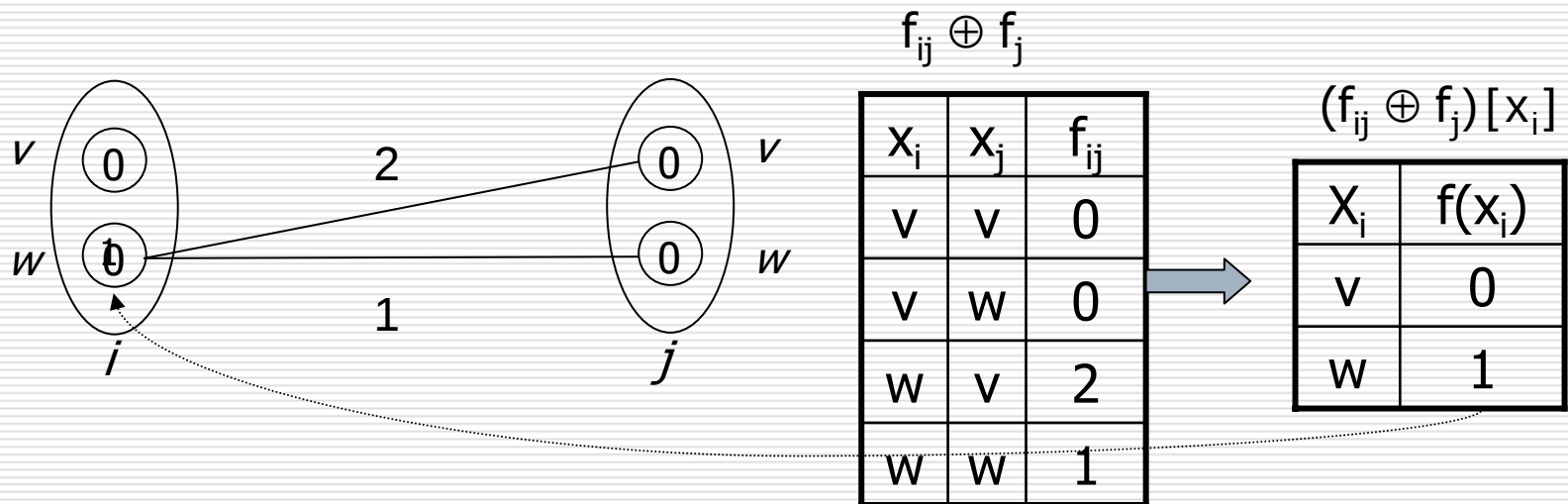
■ $c = (c_{ij} \bowtie c_j)[x_i]$ brings no new information on x_i



Arc consistency and soft constraints

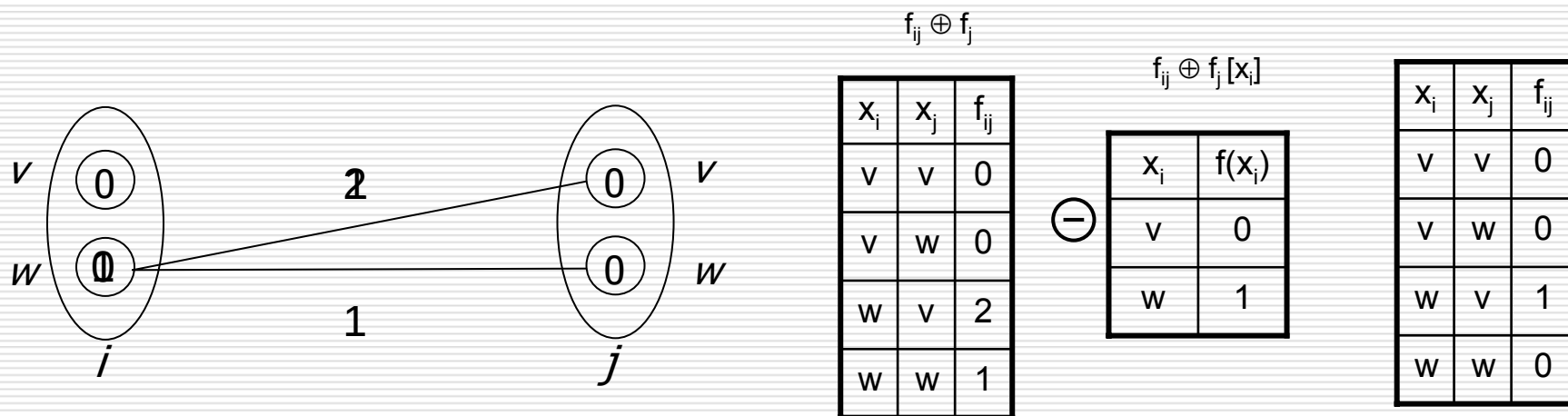
□ for any x_i and f_{ij}

■ $f = (f_{ij} \oplus f_j)[x_i]$ brings no new information on x_i



Always equivalent iff $\oplus$ idempotent [Rosenfeld 76, Bistarelli et al. 95]

Extraction of implied constraints



- Combination+Extraction: equivalence preserving transformation $[S, CS]$
- Requires the existence of a pseudo difference $(a \ominus b) \oplus b = a$: fair valued CN.

Fair valuation structures

- Totally characterized $[CS, C]$
 - Slices that interact together idempotently
 - Each slice (discrete case) is isomorphic to
 - $S_{\infty} = (\mathbb{N} \cup \{\infty\}, +)$ infinite top
 - $S_k = ([0, k], +_k)$ finite top

This means only **3 types** of structures to consider.

Tractability only known for finite costs+infinite T
(idempotent case should be...easy)

(K,Y) equivalence preserving inference

- For a set K of constraints and a scope Y
 - Replace K by $(\oplus K)$
 - Add $(\oplus K)[Y]$ to the CN (implied by K)
 - Extract $(\oplus K)[Y]$ from $(\oplus K)$

- Yields an equivalent network
- All implicit information on Y in K is explicit

- Repeat for a class of (K,Y) until fixpoint

Node Consistency (NC^{*}): $(f_i, \emptyset)$ EPI

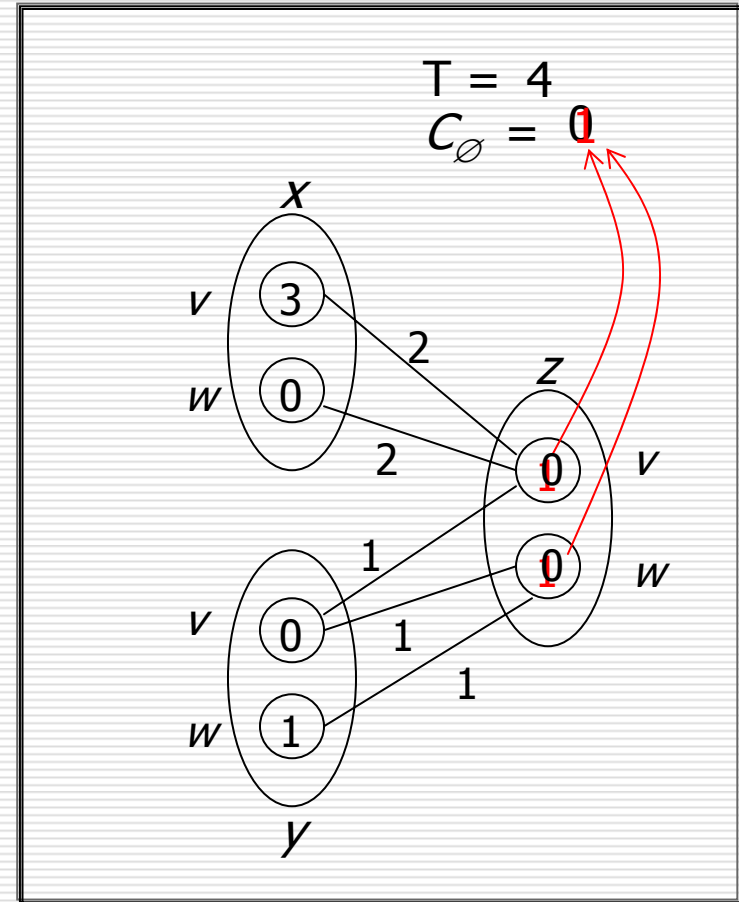
■ For any variable X_i

□ $\forall a, f_{\emptyset} + f_i(a) < T$

□ $\exists a, f_i(a) = 0$

■ Complexity:

$O(nd)$



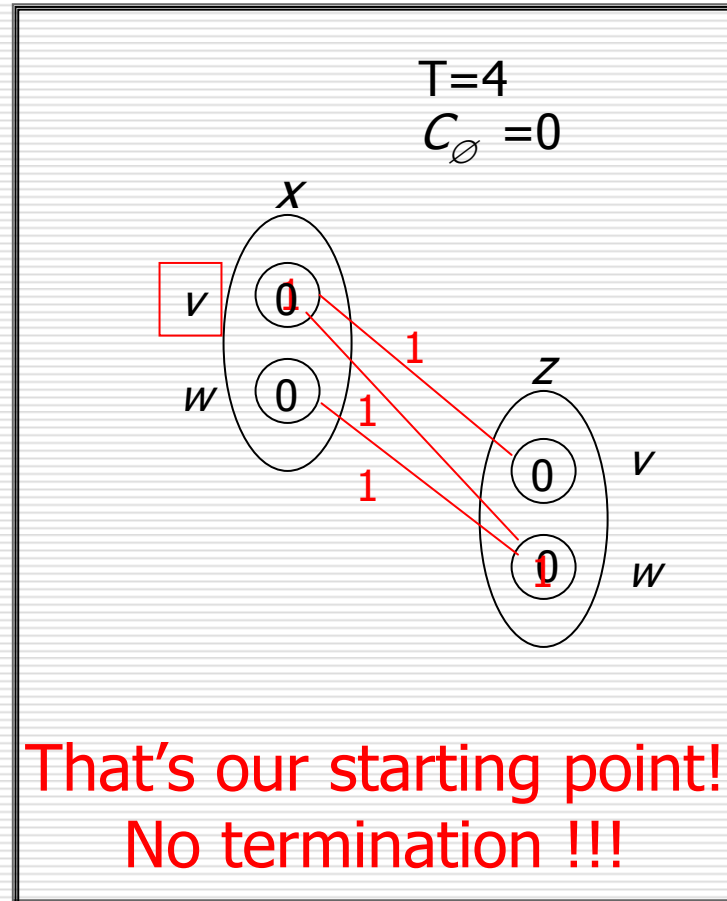
Full AC (FAC^{*}): ($\{f_{ij}, f_j\}, x_i$) EPI

■ NC^{*}

■ For all f_{ij}
 $\square \forall a \exists b$

$$f_{ij}(a, b) + f_j(b) = 0$$

(full support)



Arc Consistency (AC^*): ($\{f_{ij}\}, x_i$) EPI

■ NC^*

■ For all f_{ij}

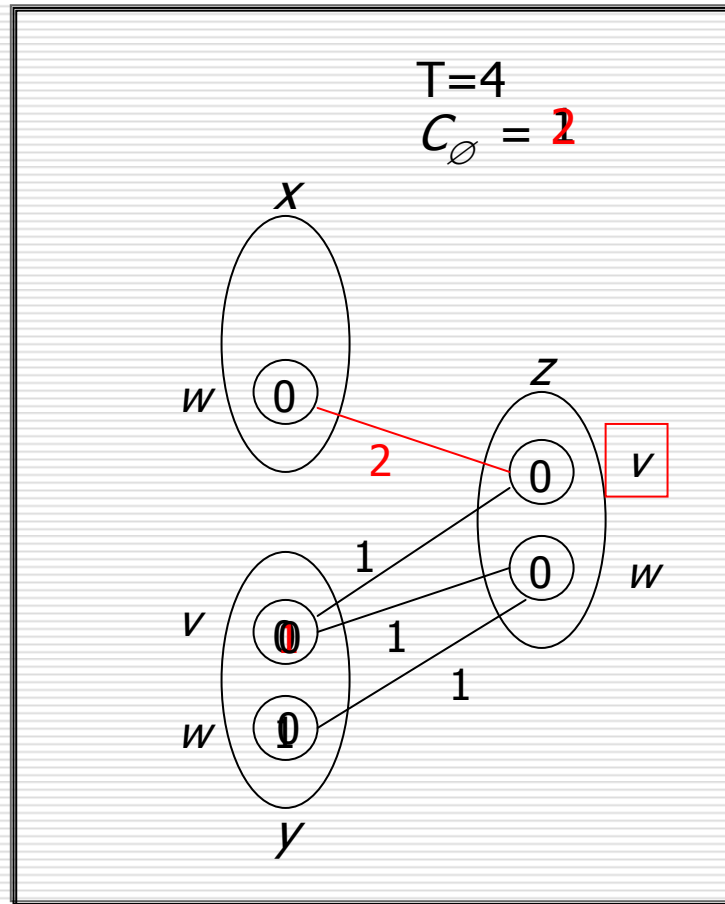
□ $\forall a \in b$

$f_{ij}(a, b) = 0$

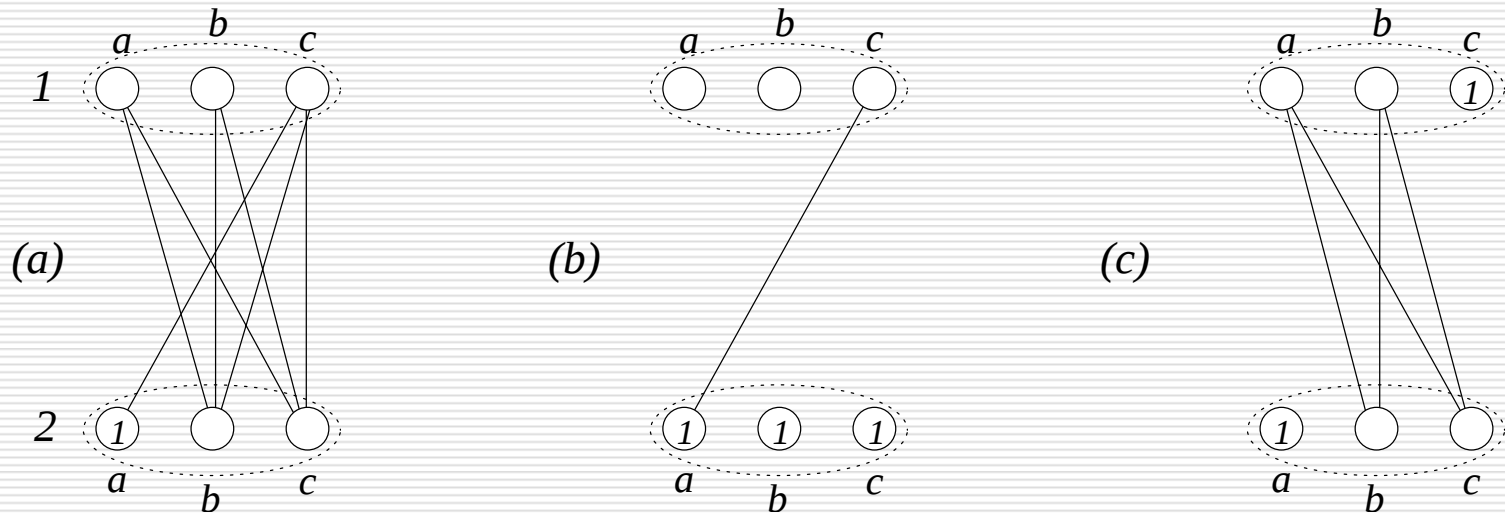
■ b is a *support*

■ complexity:

$O(n^2 d^3)$



Confluence is lost



Finding an AC closure that maximizes the lb is an NP-hard problem [CS]

Directional AC (DAC^{*}): $(\{f_{ij}, f_j\}, x_i)_{i < j}$ EPI

■ NC^{*}

■ For all f_{ij} ($i < j$) $\forall a \exists b \quad f_{ij}(a, b) + f_j(b) = 0$

■ b is a *full-support*

■ complexity: $O(ed^2)$ and it solves trees.

Simultaneous AC operations ap.

- Use rationals for costs. Infinite T.
- Value a , variable i , cost function f_{ij} .
 - $w_{iaj} = \text{cost sent from } f_i(a) \text{ to } f_{ij} - \text{from } f_{ij} \text{ to } f_i(a).$
 - $w_i = \text{amount sent from } f_i \text{ to } f_{\emptyset}.$
- $\text{Max } \sum w_i$
 - For all (a,i) $f_i(a) - \sum_j w_{iaj} - w_i \geq 0$
 - For all f_{ij}, a, b $f_{ij}(a, b) + w_{iaj} + w_{jbi} \geq 0$

Solved in pol. time by LP for bounded arities

Extra properties

- More powerful than any seq. of AC preserving transformation (infinite T)
- No better equivalent problem with the same scopes (finite costs, infinite T)
 - Quite expensive in practice.

Pb	nvar / ncost	EDAC	cpu"	LPb	cpu"	ub
CELAR06	100/1332	0	0.02	3.5	621	3389*
CELAR07	200/2865	10000	0.04	31453	3529	343592

Max-Flow as local consistency ?

If $d=2$ (Max-2SAT), the LP is a max-flow problem

- What if $d > 2$?
- What if finite T ?

Current only tractable language of MaxCSP

- 2 Monotone (maxSAT)
- Submodular (maxCSP): solved by a max-flow algorithm

Any connection ?

References

- ❑ Tutorial slides on cost function optimization (soft constraints) are available here: <http://www.math.unipd.it/~frossi/cp-school/Ecole2005.ppt>.
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- ❑ Rosenfeld, A., Hummel, R., and Zucker, S. (1976). "Scene labeling by relaxation operations". IEEE Transactions on Systems, Man and Cybernetics , 6:420--433.
- ❑ [Valued Constraint Satisfaction Problems: hard and easy problems](#) *T. Schiex, H. Fargier et G. Verfaillie* Proc. of the International joint Conference in AI, Montreal, Canada, August 1995.
- ❑ [Semiring-based CSPs and Valued CSPs: Basic Properties and Comparison](#). *S. Bistarelli, H. Fargier, U. Montanari, F. Rossi, T. Schiex, et G. Verfaillie* (LNCS 1106, Selected papers from the Workshop on Over-Constrained Systems at CP'95, reprints and background papers), pages 111-150. Springer, 1996.
- ❑ [Arc consistency for Soft Constraints](#) *T. Schiex* Proc. of CP'2000. Singapour, September 2000.
- ❑ [Arc consistency for soft constraints](#), *Martin Cooper, Thomas Schiex* *Artificial Intelligence, Volume 154 (1-2), Avril 2004*, 199-227.
- ❑ High-order consistency in valued constraint satisfaction, M. Cooper. Constraints. 10:283-305, 2005