

Exercises accompanying ‘The Modelling and Analysis of Security Protocols’

Chapter 5

Exercise 5.1 In the first version of the Yahalom protocol above, investigate the effect of moving b ’s identity outside of the encrypted component of Message 3a: adapt the Casper script to model this change, and then use FDR to analyse the protocol. \triangle

Exercise 5.2 Consider the Needham Schroeder Public Key Protocol:

Message 1. $a \rightarrow b : \{na.a\}_{PK(b)}$
 Message 2. $b \rightarrow a : \{na.nb\}_{PK(a)}$
 Message 3. $a \rightarrow b : \{nb\}_{PK(b)}$

It is assumed that each agent knows all the public keys. Analyse this protocol using Casper and FDR. \triangle

Exercise 5.3 In some protocols, nonces are considered to be *predictable*: that is, an intruder may be able to predict which nonces other agents are about to use. How can we model predictable nonces within Casper? \triangle

Exercise 5.4 Produce a Casper script to model the Yahalom-BAN protocol:

Message 1. $a \rightarrow b : na$
 Message 2. $b \rightarrow s : nb.\{a.na\}_{ServerKey(b)}$
 Message 3. $s \rightarrow a : nb.\{b.kab.na\}_{ServerKey(a)}.\{a.kab.nb\}_{ServerKey(b)}$
 Message 4. $a \rightarrow b : \{a.kab.nb\}_{ServerKey(b)}.\{nb\}_{kab}$,

for example, by adapting the script from Section 5.1. \triangle

Exercise 5.5 Adapt the script for the seven message adapted Needham Schroeder Public Key Protocol to remove the identities from within the encrypted components of the key delivery messages, messages 3 and 6:

Message 1. $a \rightarrow s : b$
 Message 2. $s \rightarrow a : \{PK(b)\}_{SSK(s)}$
 Message 3. $a \rightarrow b : a, b, \{na, a\}_{PK(b)}$
 Message 4. $b \rightarrow s : a$
 Message 5. $s \rightarrow b : \{PK(a)\}_{SSK(a)}$
 Message 6. $b \rightarrow a : b, a, \{na, nb, b\}_{PK(a)}$
 Message 7. $a \rightarrow b : a, b, \{nb\}_{PK(b)}$.

Analyse this protocol using Casper and FDR. \triangle

Exercise 5.6 Consider the following protocol, based upon one suggested by Denning and Sacco:

Message 1. $a \rightarrow b : cert_A$
 Message 2. $b \rightarrow a : cert_B.\{\{kab\}_{SK(b)}\}_{PK(a)}$

Here $cert_A$ and $cert_B$ are public key certificates for a and b , respectively, of the following form:

$$\begin{aligned} cert_A &= \{a, PK(a)\}_{CASK(ca)}, \\ cert_B &= \{b, PK(b)\}_{CASK(ca)}, \end{aligned}$$

where $CASK(ca)$ is the secret key of trusted certification authority ca .

Analyse this protocol using Casper and FDR. Hint: use environment messages to model the agents retrieving their key certificates; you should include the intruder's key certificate in his initial knowledge. \triangle

Exercise 5.7 Investigate which, if any, authentication specifications are satisfied by the Needham-Schroeder Public Key Protocol from [1]; also investigate whether the corrected version in [1] satisfies the **StrongSecret** specification for the nonces. \triangle

Exercise 5.8 Code up the Needham Schroeder Signature Protocol from Section 5.5 of the book (Hash Functions) as a Casper script. Compile the script into CSP, and analyse it using FDR. \triangle

Exercise 5.9 Consider the TMN protocol:

Message 1. $A \rightarrow S : B.\{ka\}_{pks}$
 Message 2. $S \rightarrow B : A$
 Message 3. $B \rightarrow S : A.\{kb\}_{pks}$
 Message 4. $S \rightarrow A : ka \oplus kb$

where pks is the public key of server s , ka and kb are session keys, and the intention is to establish a new session key kb shared between A and B . Use Casper and FDR to analyze this protocol; you should discover an attack. Suggest how to adapt the protocol to prevent this attack, and then investigate whether the adapted protocol is secure. \triangle

Bibliography

- [1] M. Burrows, M. Abadi, and R. Needham, *A Logic of Authentication*, Digital Equipment Corporation Systems Research Center report No. 39, 1989.